


Basics of programming


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Recursive functions

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Recursive functions

■ Problem

□ Let's calculate the n^{th} factorial!

```
int fact(int x) {
    int i, y;
    y = 1;
    for (i = 2; i <= x; i++) {
        y *= i;
    }
    return y;
}
```

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Recursive functions

■ Problem

- Let's calculate the n^{th} factorial!
- Without loops: $n! = n \cdot (n-1)!$

```
int fact(int x) {  
    if (x == 0) return 1;  
    else return x*fact(x-1);  
}
```

Recursive functions

■ Sometimes recursion is simpler

■ Problem

- calculate n^{th} fibonacci number!
- definition: $a_1 = 1, a_2 = 1, a_n = a_{n-1} + a_{n-2}$

```
int fibo(int x) {  
    if (x==1 || x==2) return 1;  
    else return fibo(x-1)+fibo(x-2);  
}
```

Recursive functions

■ Try fibonacci with loops!

- definition: $a_1 = 1, a_2 = 1, a_n = a_{n-1} + a_{n-2}$

```
int fibo(int x) {  
    int i, a1, a2, a3;  
    a1 = a2 = a3 = 1;  
    for (i = 3; i <= x; i++) {  
        a2 = a3;  
        a1 = a2;  
        a3 = a1+a2;  
    }  
    return a3;  
}
```

Towers of hanoi

■ Problem:

- we have three rods and n disks
- move the disks on first rod to the second
- only 1 piece a time and always put smaller on bigger



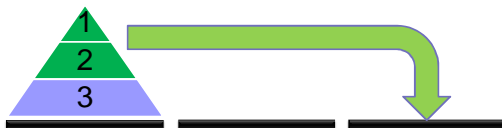
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Towers of hanoi solution

■ Print out each move

```
void hanoi(int n, int from, int to, int help) {  
    if (n == 0) return;  
    hanoi(n-1, from, help, to);  
    printf("disk %d, %d -> %d\n", n, from, to);  
    hanoi(n-1, help, to, from);  
}
```



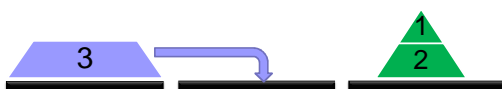
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Towers of hanoi solution

■ Print out each move

```
void hanoi(int n, int from, int to, int help) {  
    if (n == 0) return;  
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    printf("disk %d, %d -> %d\n", n, from, to);  
    hanoi(n-1, help, to, from);  
}
```



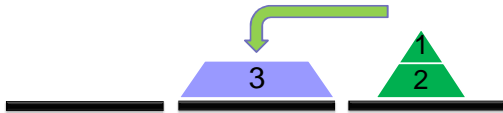
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Towers of hanoi solution

■ Print out each move

```
void hanoi(int n, int from, int to, int help) {  
    if (n == 0) return;  
    hanoi(n-1, from, help, to);  
    printf("disk %d, %d -> %d\n", n, from, to);  
    hanoi(n-1, help, to, from);  
}
```



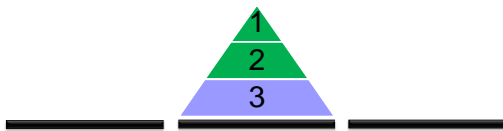
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Towers of hanoi solution

■ Print out each move

```
void hanoi(int n, int from, int to, int help) {  
    if (n == 0) return;  
    hanoi(n-1, from, help, to);  
    printf("disk %d, %d -> %d\n", n, from, to);  
    hanoi(n-1, help, to, from);  
}
```



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Rules of recursion

- Only one step is to be solved
 - e.g.: moving a single disk
- Rest is done by a new call
 - e.g.: moving n-1 disks
- Always have an exit
 - e.g.: 0 disks needs no move
- Local variables are unique in each call
- Too deep recursion should be avoided

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Exercise

- Solve factorial problem
- Print out in each function
 - the depth of recursion
 - the address of parameter x
 - use %p in *printf*: `printf("%p\n", &x);`

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Recursive data structures

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Store data in order

- Let's store data in order
 - array
 - with each insertion $n/2$ elements should be moved on average
 - binary tree
 - stores elements in nodes of a tree
 - each element is usually at most $\log_2 n$ far from root

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Binary tree

■ Data structure

```
typedef struct btree {  
    int n;  
    struct btree *left, *right;  
} btree;
```

■ Rules

□ leaves have null pointers for children

□ for each node x

■ $x \rightarrow \text{left} \rightarrow n < x \rightarrow n$

■ $x \rightarrow \text{right} \rightarrow n \geq x \rightarrow n$

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Binary tree: empty tree and creation

■ An empty tree is a NULL pointer

```
btree* tree = NULL;
```

■ Create one element tree

```
btree* create(int n) {  
    bintree* t;  
    t = (btree*)malloc(sizeof(btree));  
    t->left = t->right = NULL;  
    t->n = n;  
    return t;  
}
```

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Binary tree: insertion

■ Algorithm

■ Insert number n into tree t

if t is empty

n is new tree

otherwise

if $n < (t \rightarrow n)$, insert into $t \rightarrow \text{left}$

otherwise insert into $t \rightarrow \text{right}$

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Binary tree: insertion

- Insert number n into tree t

```
btree* insert(btree* t, int n) {
    if (t == NULL) {
        t = create(n);
    } else {
        if (n < t->n)
            t->left = insert(t->left, n);
        else
            t->right = insert(t->right, n);
    }
    return t;
}
```

Binary tree: contains

- Contains: return true if *tree* contains n

```
int contains(btree* t, int n) {
    if (t == NULL) {
        return 0;
    } else {
        if (t->n == n)
            return 1;
        else if (n < t->n)
            return contains(t->left, n);
        else
            return contains(t->right, n);
    }
}
```

Binary tree traversal

- How can we traverse a binary tree?

- ☐ inorder
 - left subtree, root, right subtree
 - e.g. printing data in increasing order
- ☐ preorder
 - root, left subtree, right subtree
- ☐ postorder
 - left subtree, right subtree, root

Binary tree algorithms

- Most algorithms are recursive
 - ☐ empty tree
 - ☐ root node
 - ☐ left or right subtree
- Algorithms for exercise
 - ☐ depth of tree
 - ☐ number of nodes, leaves, parents
 - ☐ printing out: inorder, preorder, postorder
 - ☐ deleting a tree

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